Unitary non-backtracking eigenvalues



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Leo Torres

660

PhD in Network Science

Network Science Institute, Northeastern University

leo@leotrs.com





G = (V, E)



• community detection

- Krzakala, et al. PNAS 110.52 (2013): 20935-20940.
- Bordenave, et al. FOCS (2015).

• centrality

- Martin, et al. Phys. Rev. E 90.5 (2014): 052808.
- Morone & Makse. Nature 524.7563 (2015): 65-68.
- Arrigo, et al. J. of Sci. Comp. 80.3 (2019): 1419-1437.

• dynamics (SIR, SIS)

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- Hamilton, & Pryadko. Phys. Rev. Lett. 113.20 (2014): 208701.
- Shrestha, et al. Phys. Rev. E 92.2 (2015): 022821.
- Castellano, & Pastor-Satorras. Phys. Rev. E 98.5 (2018): 052313.
- Masuda, et al. J. of App. Math. 85.2 (2020): 214-230.

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most graphs I've seen have a simple spectrum...

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Questions

- What's different about the unitary eigenvalues?
- Can we compute their multiplicities?
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Let G be a graph with non-backtracking matrix B. Let $B\mathbf{v}=\lambda\mathbf{v}$ and $|\lambda|=1.$

Tasks:

- study \mathbf{v}, λ
- compute the number of such **v** that are L.I.



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Recall:

- B has one row/column for each oriented edge (size 2m imes 2m)
- **v** is a function of the oriented edges

• As a matrix:

$$B_{k
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• As an operator: "v into k" $(B\mathbf{v})_{k
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What's more: it can be shown that nodes of degree one do not influence the non-zero eigenvalues.





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From now on, assume the graph has minimum degree 2.



• If **v** is an eigenvector:

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More assumptions...

From now on:

- 1. finite
- 2. simple
- 3. undirected
- 4. unweighted
- 5. connected
- 6. minimum degree 2
- 7. not a circle graph (i.e. at least one node w degree 3+)









 $\lambda^2=1$




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Graph subdivisions

Definition: For a graph *G*, its *p*-th subdivision is the graph formed by replacing each edge by a chain of length *p*.



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This is what's special about unitary e'values.

Theorem: Let H be the p-th subdivision of some graph G s.t. G is not the subdivision of any other graph. Suppose λ is a p-th root of unity. Then $GM(\lambda) = |E| - |V| + 1$.

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Sketch:

1. Define the zeta function

$$Z_G(t) = \prod_{c \in G} \left(1 - t^{|c|}
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- Step 2: build e'vectors of *H* using e'vectors of *G* (and vice-versa).

• Step 3:
$$GM_H(\lambda) = GM_G(1)$$
.
 $\begin{cases} ec{\mathbf{v}}^l = 0 & ext{for each node } l, \\ \mathbf{v}_{k o l} + \mathbf{v}_{l o k} = 0 & ext{for each edge } k - l. \end{cases}$ $|E| \stackrel{+1}{-1} \stackrel{-1}{-1} + 1$

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Some notation

Definition: True nodes and subdivision nodes.



 S_p

Identify some of the true nodes of S_p with a node in G.



G

Identify some of the true nodes of S_p with a node in G.



 S_p



H

G

Identify some of the true nodes of S_p with a node in G.

Lemma: Take an eigenvector of S_p . Pad it with zeros. Then that vector is "non-leaky" in H, i.e. $(d_l - 2) \vec{\mathbf{v}}^l = 0, \forall l$.



Some notation

Definition: if the subdivision is even, we also have middle nodes.



Identify some of the true or middle nodes nodes of S_p with a node in G.

Lemma: Take an eigenvector of S_p . Pad it with zeros. Then that vector is "non-leaky" in H, i.e. $(d_l - 2) \vec{\mathbf{v}}^l = 0, \forall l$.





H
















G

 S_p

H

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PB